

Grover's algorithm
 qubitization w. basis change
 amplitude amplification (AA)
 oblivious amplitude amplification

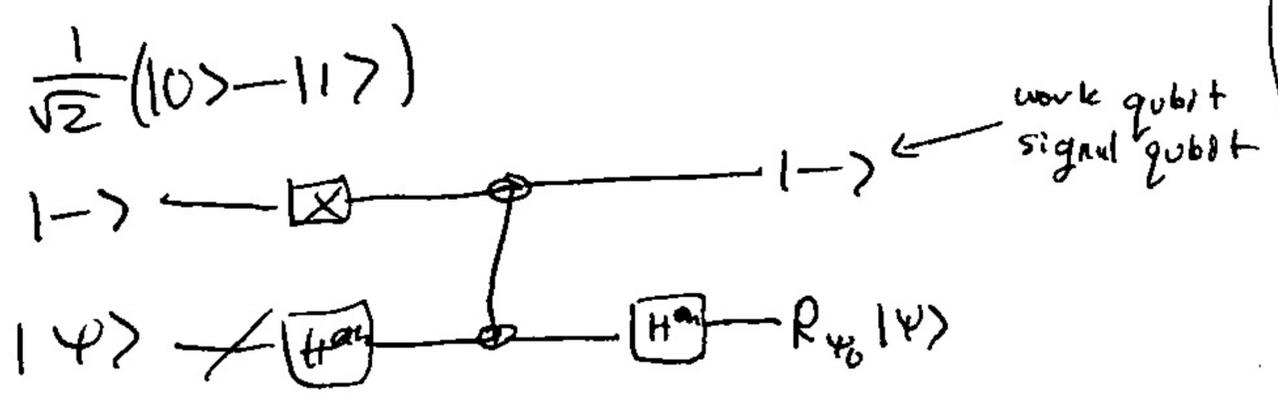
$\square \quad \boxed{3} \quad \dots \quad \square$ ^{N-boxes}
 $f: \{0,1\}^n \rightarrow \{0,1\}$ (where 1 is marked)
 $f(x) \quad R_{x_0} = I - 2|x_0\rangle\langle x_0|$
 phase kickback

$$R_{x_0}|x\rangle = (-1)^{f(x)}|x\rangle = \begin{cases} -|x_0\rangle & x=x_0 \\ |x\rangle & \text{other} \end{cases}$$

How many times do we need to query R_{x_0} (quantum unboxing)

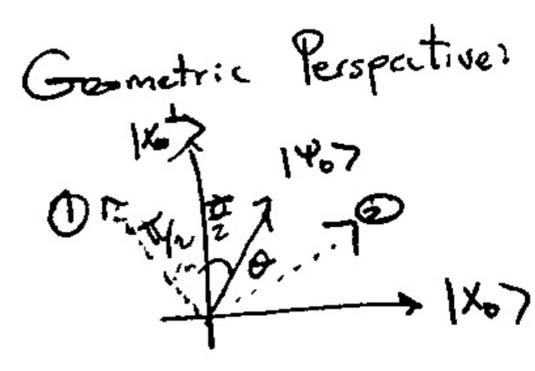
$$|\psi_0\rangle = \frac{1}{\sqrt{2^n}} \sum_{i=0}^{2^n-1} |i\rangle = H^{\otimes n} |0^n\rangle$$

$$R_{\psi_0} = 2|\psi_0\rangle\langle\psi_0| - I = H^{\otimes n} (2|0^n\rangle\langle 0^n| - I) H^{\otimes n}$$



Grover Iterate: $G = R_{\psi_0} R_{x_0}$
 claim $G^k |\psi_0\rangle$ $k \approx \Theta(\sqrt{N})$
 measure in computational basis

then w. $\Omega(1)$ prob find $|x_0\rangle$



angle $\frac{\theta}{2} \rightarrow \frac{3\theta}{2}$
 $2\theta + \frac{\theta}{2} = \frac{5\theta}{2}$

$$|x_0^\perp\rangle = \frac{1}{\sqrt{N-1}} \sum_{x \neq x_0} |x\rangle \quad |\psi_0\rangle = \frac{1}{\sqrt{N}} \sum |x\rangle$$

$$|\psi_0\rangle = \frac{1}{\sqrt{N}} |x_0\rangle + \sqrt{\frac{N-1}{N}} |x_0^\perp\rangle = \sin\frac{\theta}{2} |x_0\rangle + \cos\frac{\theta}{2} |x_0^\perp\rangle$$

$$\mathcal{B} = \{|x_0\rangle, |x_0^\perp\rangle\}$$

$$G^k |\psi_0\rangle = \sin\left(\frac{2k+1}{2}\theta\right) |x_0\rangle + \dots$$

which k ? $\frac{2k+1}{2}\theta \approx \frac{\pi}{2} \implies k \approx \left(\frac{\pi}{2\theta} - 1\right) \cdot \frac{1}{2}$
 $\left(\theta \approx \frac{2}{\sqrt{N}}\right) \approx \frac{\pi\sqrt{N}}{4}$

Grover's algorithm: "doesn't converge"

M mark vertices $f(x) = \begin{cases} 1, & x \in \text{marked set} \\ 0, & \text{else} \end{cases}$

qubitization $U_A^\dagger Z_\pi U_A Z_\pi$

$A \in \mathbb{C}^{M \times N}$ projector Π projector Π'

$$B = \{ |\psi_0\rangle, \dots, |\psi_{N-1}\rangle, \dots \} \quad B' = \{ |\psi_0\rangle, \dots, |\psi_{N-1}\rangle, \dots \}$$

$M \times N$

$$\Pi' U_A \Pi = \sum_{i,j \in [N]} |\psi_i\rangle A_{ij} \langle \psi_j|$$

$$U_A = \begin{bmatrix} A & * \\ * & * \end{bmatrix}$$

"more general version of block encoding"

deficient ~~$\tilde{U}_A = [U_A]_{B'}^{B'}$~~
 coefficient $\tilde{U}_A = [U_A]_{B'}^{B'}$

$$\tilde{U}_A^\dagger Z_{\pi_0 m} \tilde{U}_A Z_{\pi_0 n} \quad \text{--- [lec 10.6]}$$

Interpret Grover's alg using qubitization

$$B = \{ |\psi_0\rangle, \dots \} \quad B' = \{ |x_0\rangle, \dots \}$$

"odd order polynomial"

$$U_A = R_{\psi_0}$$

$$[U_A]_{B'}^{B'} = \begin{bmatrix} a & * \\ * & * \end{bmatrix}$$

$$a = \sin \frac{\theta}{2} = \frac{1}{\sqrt{N}}$$

$$Z_\pi = R_{\psi_0}$$

$$Z_{\pi'} = -R_{x_0}$$

$$U_A Z_\pi (U_A^\dagger Z_{\pi'} U_A Z_\pi)^k$$

$$U_A Z_\pi = R_{\psi_0} R_{\psi_0} = I$$

$$(U_A^\dagger Z_{\pi'})^k = (-1)^k (R_{\psi_0} R_{x_0})^k$$

$$\rightarrow \begin{bmatrix} (-1)^k T_{2k+1}(a) & * \\ * & * \end{bmatrix}$$

$$\sin \left(\frac{(2k+1)\theta}{2} \right) |x_0\rangle + | \perp \rangle$$

$$T_{2k+1}(a) = (-1)^k \sin((2k+1) \arcsin a)$$

$$* |\psi_0\rangle = T_{2k+1}(a) |x_0\rangle \langle \psi_0|$$

"AA and AE paper"

$$U_{\psi_0} |0^m\rangle = |\psi_0\rangle$$

$$|\psi_0\rangle = \sqrt{p} |\psi_{\text{good}}\rangle + \sqrt{1-p} |\psi_{\text{bad}}\rangle$$

solution $|\psi_{\text{good}}\rangle$ w. (0) + $\mathcal{O}(\frac{1}{\sqrt{p}})$

$$R_{\text{good}} = I - 2|\psi_{\text{good}}\rangle\langle\psi_{\text{good}}|$$

$$G = R_{\psi_0} R_{\text{good}}$$

$$G^k \quad k \sim \frac{\pi}{4\sqrt{p}}$$

$$|\psi_{\text{good}}\rangle = |0^m\rangle |\psi_{\text{good}}\rangle$$

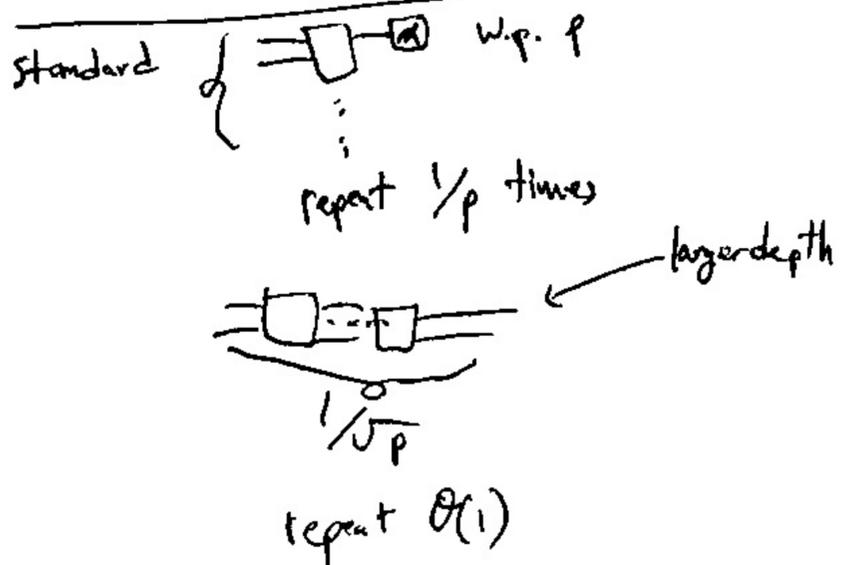
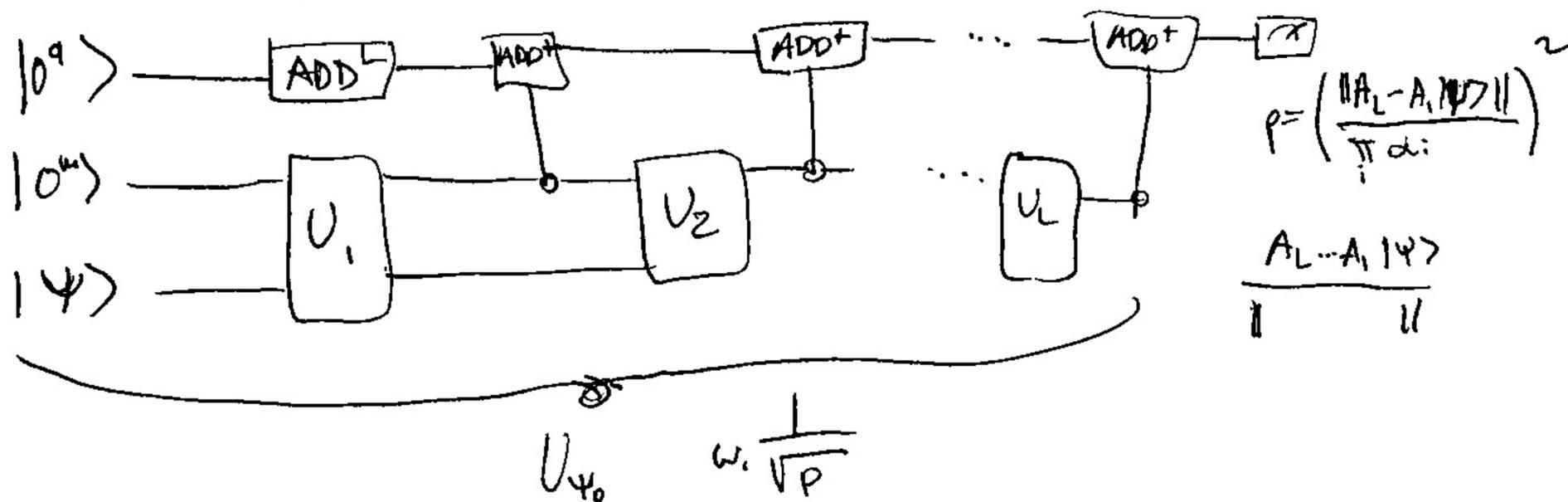
Revisit composition of Block Encodings

$U_1, \dots, U_L \quad U_i \in \text{BE}_{a_i, m}(A_i)$ goal: block encoding of $A_L \dots A_1$

Naive way: use $m \cdot L$ qubits subnormalization $\prod \alpha_i$

more efficient: L additional qubits subnormalization $\prod \alpha_i$

even more efficient: $\lceil \log L \rceil$



"AA and AE paper"

$$U_{\psi_0} |0^n\rangle = |\psi_0\rangle$$

$$|\psi_0\rangle = \sqrt{p} |\psi_{\text{good}}\rangle + \sqrt{1-p} |\psi_{\text{bad}}\rangle$$

solution $|\psi_{\text{good}}\rangle$ w. $(0) + \mathcal{O}(\frac{1}{\sqrt{p}})$

$$R_{\text{good}} = I - 2|\psi_{\text{good}}\rangle\langle\psi_{\text{good}}|$$

$$G = R_{\psi_0} R_{\text{good}}$$

$$G^k \quad k \sim \frac{\pi}{4\sqrt{p}}$$

$$|\psi_{\text{good}}\rangle = |0^n\rangle |\psi_{\text{good}}\rangle$$

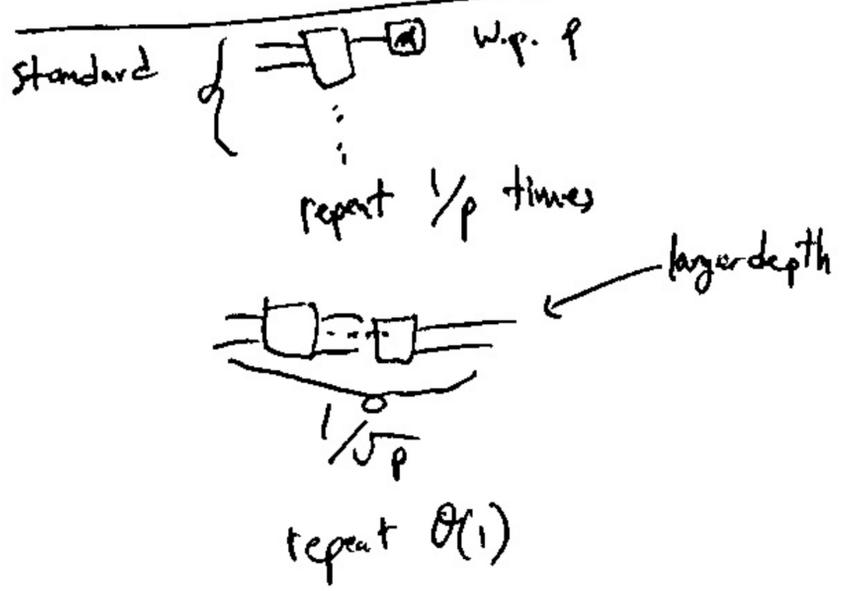
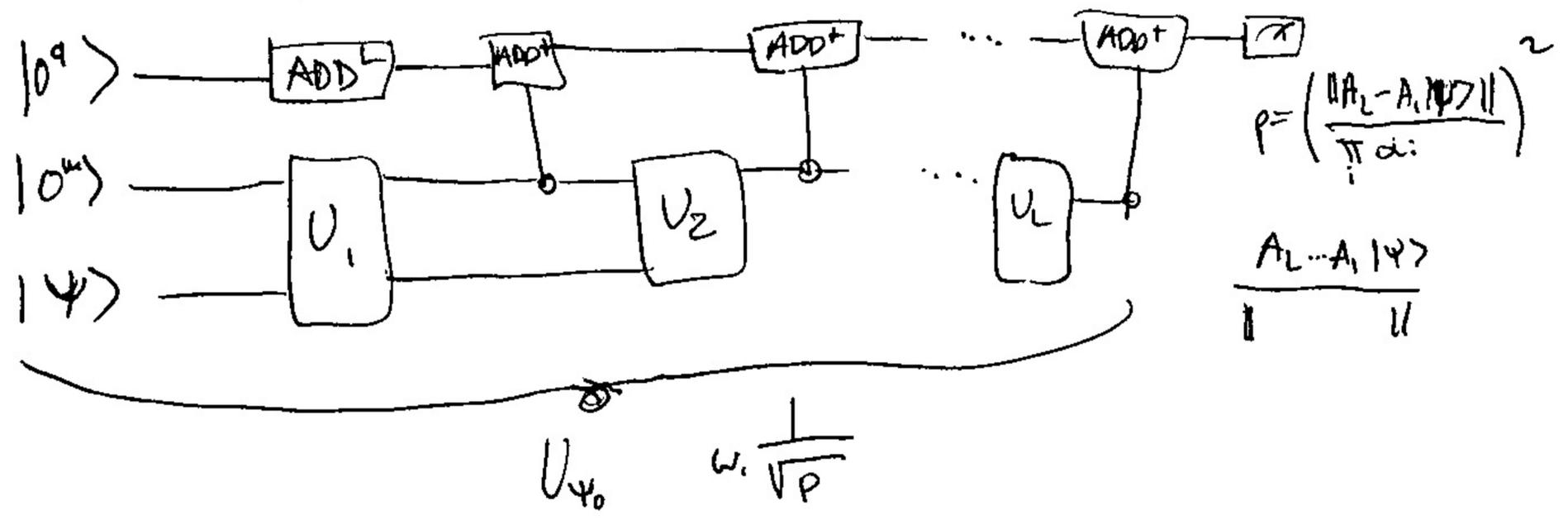
Revisit composition of Block Encodings

$U_1, \dots, U_L \quad U_i \in \text{BE}_{a_i, m}(A_i)$ goal: block encoding of A_L, \dots, A_1

Naive way: use $m \cdot L$ qubits subnormalization $\prod \alpha_i$

more efficient: L additional qubits subnormalization $\prod \alpha_i$

even more efficient: $\lceil \log L \rceil$



oblivious AA

special case: block encoding of a unitary matrix

$$e^{-iHt}$$

$$V = \begin{pmatrix} \frac{1}{\alpha} U & * \\ * & * \end{pmatrix}$$

$$V \in BE_{d,m}(U)$$

$$V|0^m\rangle|\psi\rangle = \frac{1}{\alpha}|0^m\rangle U|\psi\rangle + \dots$$

$$A = \frac{U}{\alpha} = U \cdot \frac{1}{\alpha} \cdot I$$

$$T_{2k+1}^\diamond(A)$$

~~$$T_{2k}^\diamond(A)$$~~

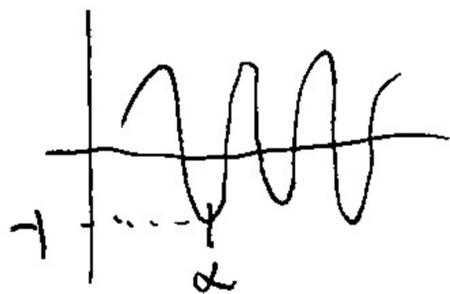
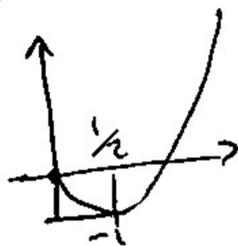
||

$$T_{2k+1}(\frac{1}{\alpha}I)$$

||

|

$$T_3(x) = 4x^3 - 3x$$



"quantum circuit"
"i.b.l"